

Date: Tue, 04 Apr 2000 18:14:23 -0400 (EDT) From: Graham Solomon  
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Subject: fulfillability  
To: Allen Hazen <a.hazen@philosophy.unimelb.edu.au> Reply-to: Graham  
Solomon <gsolomon@wlu.ca> MIME-version: 1.0

Dear Allen, I had a bit of spare time ...

Notes taken by William Demopoulos from a lecture by Hilary Putnam,  
July 21 1983, on Kripke's concept of fulfillability.

Notation:

$s = \text{finite } \{s(1), s(2), \dots, s(n)\}$  or infinite  $\{s(1), s(2), \dots, s(n), \dots\}$  of natural numbers.

$s(i)$  is the  $i$ -th term of  $s$

$l(s) = n$  or  $\infty$ , as appropriate, is the length of  $s$

{GS: I'll use  $A$  for the universal quantifier,  $E$  for the existential}

Illustration of our basic concept:  $s$  fulfills the sentence  $\phi$

let  $\phi = \forall x \exists y \forall z. \psi(x, y, z)$  where  $\psi$  is quantifier-free.

let  $s$  be a sequence of length  $v$

Then we say  $s$  fulfills  $\phi$  if  $\forall i < v$  [ i.e., for all  $i$  less than  $v$  ]  
the sentence  $\phi'$  holds in  $N$  (= the standard model of PA) where  
 $\phi' = (\forall x < s(i)) (\exists y < s(i+1)) (\forall z < s(j)). \psi(x, y, z)$  where  $i$  is  
less than or equal  $j$  less than  $v$

Note: were  $\phi$  quantifier-free,  $s$  fulfills  $\phi$  iff  $s$  satisfies  $\phi$   
(since  $\phi = \phi'$  in that case)

4 FACTS regarding fulfills:

FACT 1  $\phi$  is true iff there is an infinite increasing sequence  $s$   
which fulfills  $\phi$ .

( $\Rightarrow$ ) if  $\phi$  is true, every sequence satisfies  $\phi$ . Thus choose an  
increasing sequence with the appropriate properties

( $\Leftarrow$ ) if  $s$  is infinite and increasing such that  $s$  fulfills  $\phi$  then  
 $(\forall x \in \mathbb{N})(\exists y \in \mathbb{N})(\forall z \in \mathbb{N}) \mathbb{N} \models \psi(x,y,z)$   
 But then  $\mathbb{N} \models \phi$

Definition: A sequence  $s$  is *\_good\_* if

- (i)  $s$  is finite  $[(\exists n)(l(s) = n)]$
- (ii)  $s(i)^2 \leq s(i+1)$  for  $i < l(s)$
- (iii)  $l(s) < s(1)$  [ $s$  is "large"]

Remark: (ii) ensures that  $s$  is closed under plus and times, in the sense that  $s$  fulfills  $(\forall x,y)(\exists z)(x+y=z)$  and  $((\forall x,y)(\exists z)(x \cdot y=z))$

Definition: A sentence  $\phi$  is *\_n-fulfilled\_* if there is a good sequence  $s$  such that  $l(s) = n$  and  $s$  fulfills  $\phi$

Formally:  $(\exists s)(s \text{ is good} \ \& \ l(s) = n \ \& \ s \text{ fulfills } \phi)$

Definition:  $\phi$  is *\_fulfilled\_* if  $(\forall n)(\phi \text{ is } n\text{-fulfilled})$

FACT 2  $\phi$  is true  $\implies$   $\phi$  is fulfilled

If  $\mathbb{N} \models \phi$  then  $(\forall n)(\exists s)(l(s) = n \ \& \ \dots \ \& \ s \text{ fulfills } \phi)$ . Let  $s$  be an infinite increasing sequence which fulfills  $\phi$  (by FACT 1) such that  $s$  satisfies condition (ii) of good sequences. The problem is to find finite sequences which are good and which fulfill  $\phi$ , one for each  $n$ . The key difficulty is to ensure that  $s(1) > l(s)$  for each finite  $s$  we choose. We proceed as follows:

if  $s(1) > n$  we let  $t = s \uparrow_{n+1}$  {where  $\uparrow$  is the half-arrow pointing up}  
 otherwise define  $t$  as follows:

$t(1) = s(p)$

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$t(k+1) = s(p+k)$

where  $p$  is the first place where the value of  $s$  exceeds  $n$ . (We know that there must be such a  $p$  since  $s$  is infinite and increasing.)

Clearly  $t$  fulfills  $\phi$  if  $s$  does.

FACT 3 FACT 2 is provable in PA. I.e., for each sentence  $\phi$  PA  $\vdash$   $\phi \rightarrow (\forall n)(\exists s)(l(s)=n \ \& \ \dots \ \& \ s \text{ fulfills } \phi)$

I.e.,  $\text{PA} \vdash \phi \rightarrow (\forall n)(\exists s)(l(s)=n \ \& \ \dots \ \& \ (\forall x < s(i))(\exists y < s(i+1))(\forall z < s(j))\psi(x,y,z)(\exists i,j \text{ such that } i \text{ less than } n \ \& \ j \text{ less than } n))$

There is no difficulty with the quantification over finite sequences since such sequences can be coded by single numbers, thus eliminating the need for such quantification. So let  $s$  stand ambiguously for itself and for its Gödel-number.

FACT 4 Facts 1-3 hold for the joint fulfillability of a finite number of sentences. I.e.,  $\phi$  is true and  $\psi$  is true iff there is a sequence (infinite and increasing) which fulfills  $\phi$  and which fulfills  $\psi$ . If  $\phi$  is true and  $\psi$  is true then  $(\forall n)(\exists s)(\phi \text{ is } n\text{-fulfilled} \ \& \ \psi \text{ is } n\text{-fulfilled})$

THEOREM: Let  $T$  be a theory such that  
 $T$  is or contains PA  
 $T + \text{PA}$  is  $\omega$ -consistent  
 $T$  has an RE set of axioms {why RE, why not R?- because bounded quantifiers can define only up to RE?}

then ' $T$  is fulfilled' is true in  $\mathbb{N}$  but PA not  $\vdash$  ' $T$  is fulfilled'

(thus ' $T$  is fulfilled' is true in some other (non-standard) model of PA)

Cor. PA is not finitely axiomatizable

Proof of Cor. Let  $T$  finitely axiomatize PA  
 Since  $T$  finitely axiomatizes PA,  $\text{PA} \vdash \phi_T$ , where  $\phi_T$  is the conjunction of the axioms of  $T$ .  
 Hence:  $\text{PA} \vdash \phi_T \rightarrow$  ' $\phi_T$ ' is fulfilled, by FACT 3  
 $\text{PA} \vdash \phi_T$   
 $\text{PA} \vdash$  ' $\phi_T$ ' is fulfilled  
 I.e.,  $\text{PA} \vdash T$  is fulfilled, contradicting THEOREM

{next message will contain proof of THEOREM, the remaining half of these notes}